CSCI-1680 - Computer Networks

Physical Layer
Link Layer I

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Based partly on lecture notes by David Mazières, Phil Levis, John Jannotti, Peterson & Davie, Rodrigo Fonseca

Administrivia

- Snowcast milestone today!

Today

- **Physical Layer**
  - Modulation and Channel Capacity
  - Encoding
- **Link Layer I**
  - Framing

Layers, Services, Protocols

Application	Service: user-facing application. Application-defined messages
Transport	Service: multiplexing applications Reliable byte stream to other node (TCP), Unreliable datagram (UDP)
Network	Service: move packets to any other node in the network IP: Unreliable, best-effort service model
Link	Service: move frames to other node across link. May add reliability, medium access control
Physical	Service: move bits to other node across link

Physical Layer (Layer 1)

- **Responsible for specifying the physical medium**
  - Type of cable, fiber, wireless frequency
- **Responsible for specifying the signal (modulation)**
  - Transmitter varies *something (*amplitude, frequency, phase)
  - Receiver samples, recovers signal
- **Responsible for specifying the bits (encoding)**
  - Bits above physical layer -> *chips*

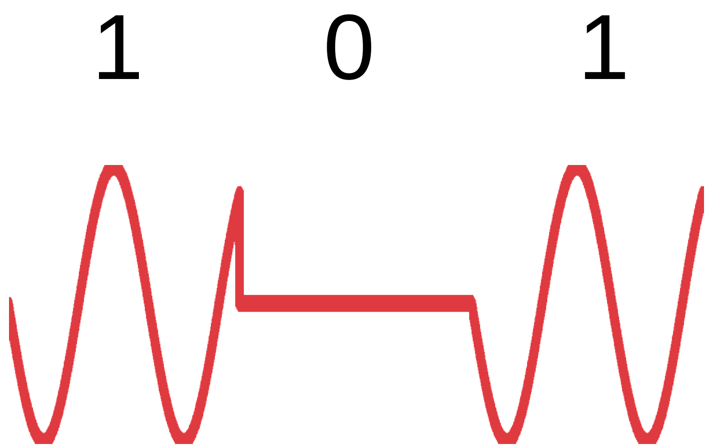
Modulation

- **Specifies mapping between digital signal and some variation in analog signal**
- **Why not just a square wave (1v=1; 0v=0)?**
  - Not square when bandwidth limited
- **Bandwidth – frequencies that a channel propagates well**
  - Signals consist of many frequency components
  - Attenuation and delay frequency-dependent

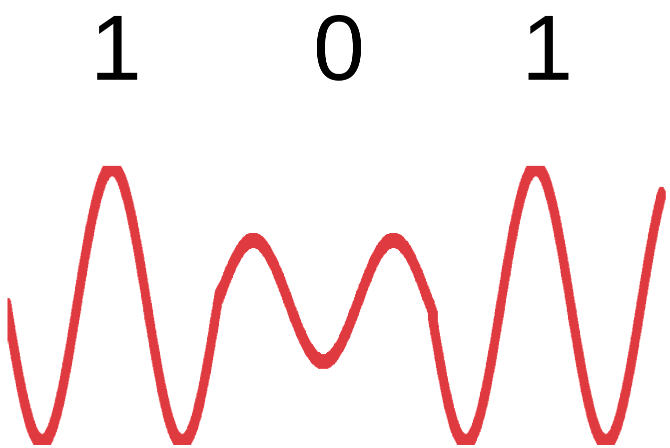
Use Carriers

- **Idea: only use frequencies that transmit well**
- ***Modulate* the signal to encode bits**

OOK: On-Off Keying



ASK: Amplitude Shift Keying

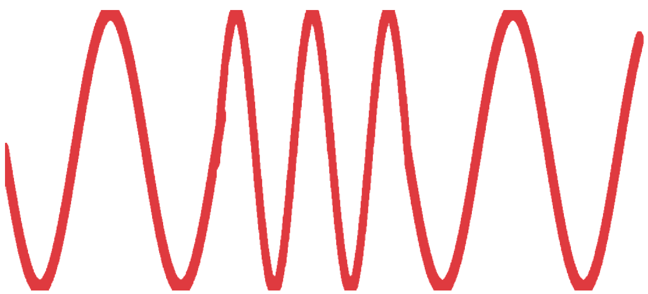


Use Carriers

- **Idea: only use frequencies that transmit well**
- ***Modulate* the signal to encode bits**

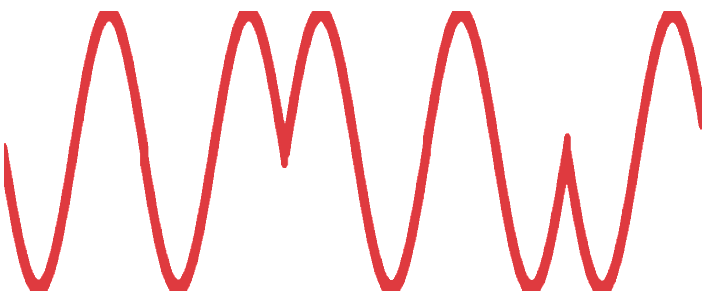
FSK: Frequency Shift Keying

1 0 1



PSK: Phase Shift Keying

1 0 1



How Fast Can You Send?

- **Encode information in some varying characteristic of the signal.**
- **If B is the maximum frequency of the signal**

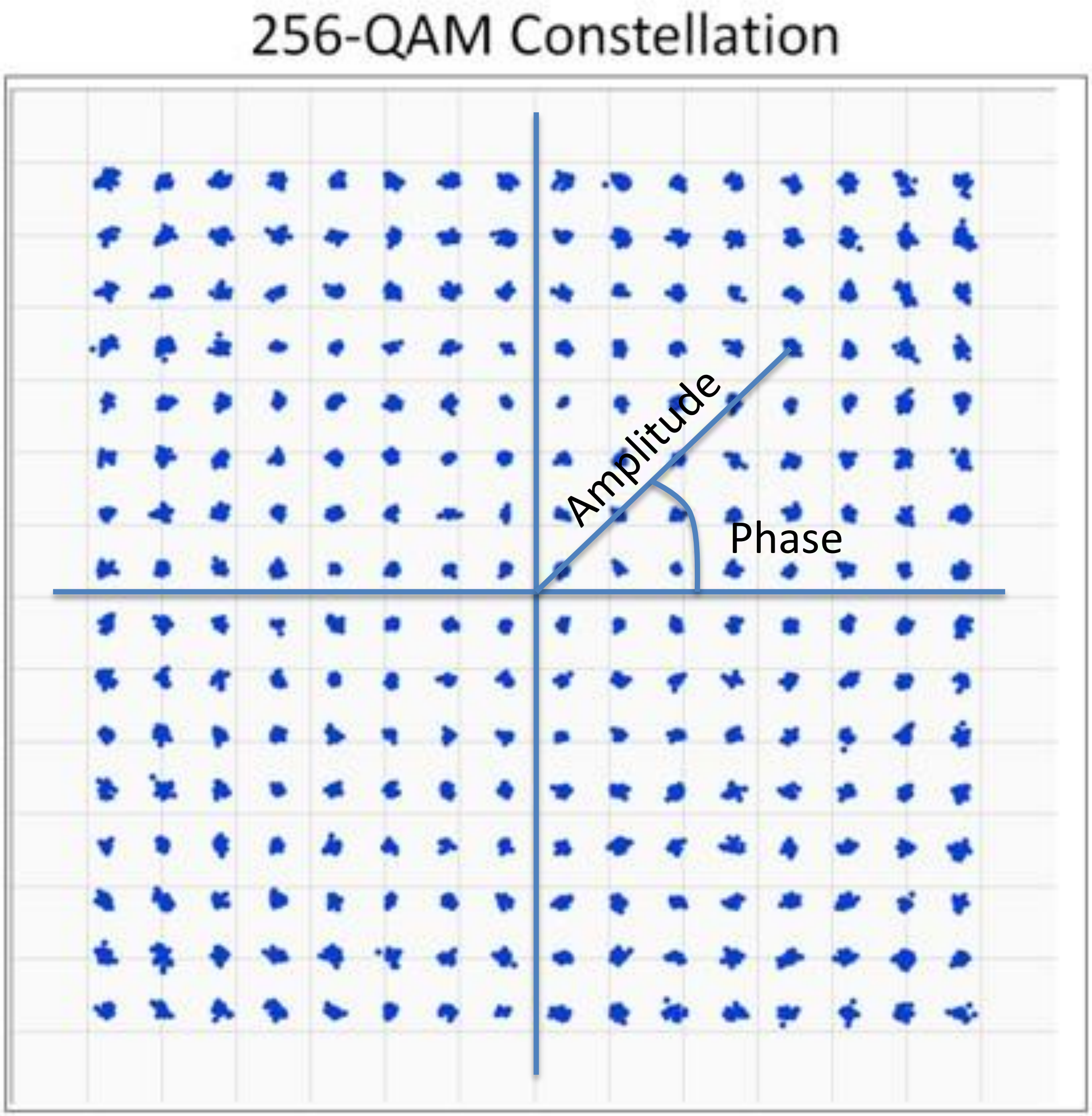
$$C = 2B \text{ bits/s}$$

(Nyquist, 1928)

Can we do better?

- **So we can only change 2B/second, what if we encode more bits per sample?**
  - Baud is the frequency of changes to the physical channel
  - Not the same thing as bits!
- **Suppose channel passes 1KHz to 2KHz**
  - 1 bit per sample: alternate between 1KHz and 2KHz
  - 2 bits per sample: send one of 1, 1.33, 1.66, or 2KHz
  - Or send at different amplitudes: A/4, A/2, 3A/4, A
  - n bits: choose among 2^n frequencies!
- **What is the capacity if you can distinguish M levels?**

Example



Hartley's Law

$$C = 2B \log_2(M) \text{ bits/s}$$

Great. By increasing M, we can have as large a capacity as we want!

Or can we?

The channel is noisy!

The channel is noisy!

- **Noise prevents you from increasing M arbitrarily!**
- **This depends on the signal/noise ratio (S/N)**
- **Shannon:** $C = B \log_2(1 + S/N)$
  - C is the channel capacity in bits/second
  - B is the bandwidth of the channel in Hz
  - S and N are average signal and noise power
  - Signal-to-noise ratio is measured in dB = $10\log_{10}(S/N)$

Putting it all together

- **Noise limits M!**

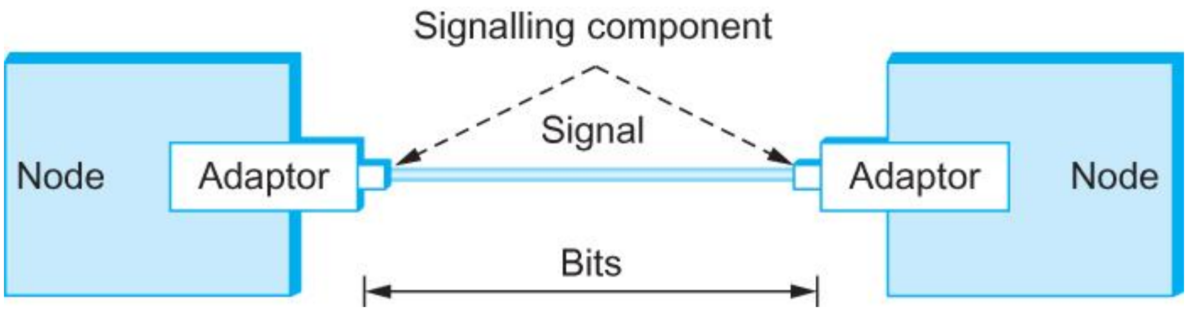
$$2B \log_2(M) \leq B \log_2(1 + S/N)$$

$$M \leq \sqrt{1+S/N}$$

- **Example: Telephone Line**
  - 3KHz b/w, 30dB S/N = 10ˆ(30/10) = 1000
  - C = 3KHz $\log_2(1001) \approx$ 30Kbps

- **Note: Until now no retransmission are allowed**

Encoding

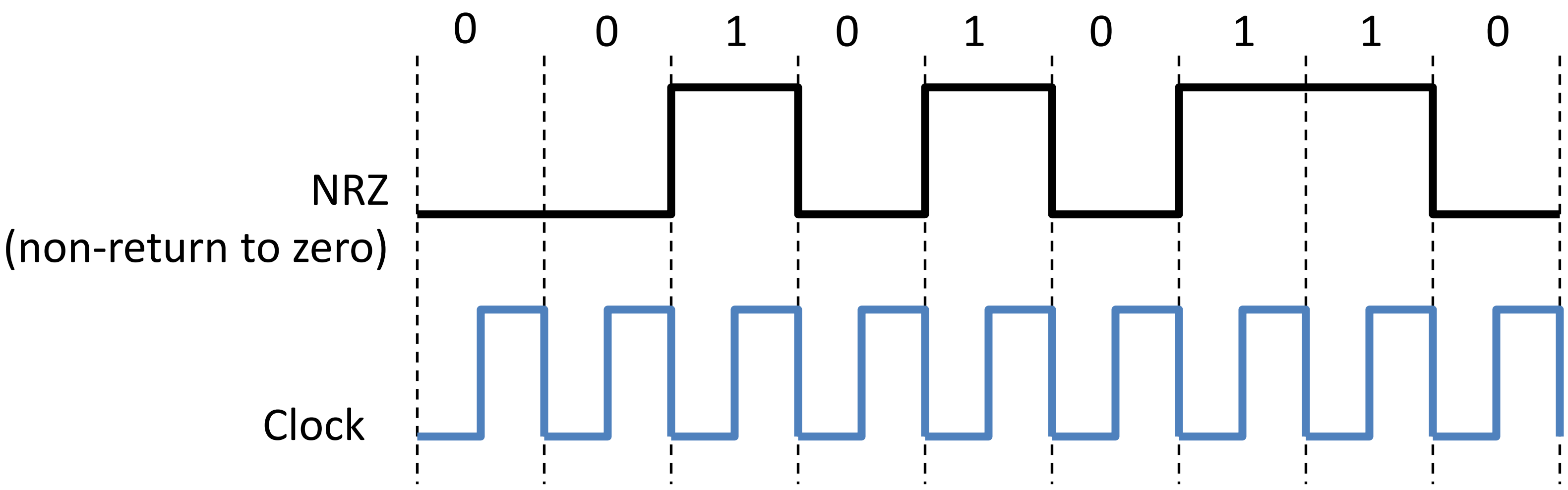


Signals travel between signaling components; bits flow between adaptors

- **Now assume that we can somehow modulate a signal: receiver can decode our binary stream**
- **How do we encode binary data onto signals?**

Encoding

- **One approach: 1 as high, 0 as low!**
  – Called Non-return to Zero (NRZ)

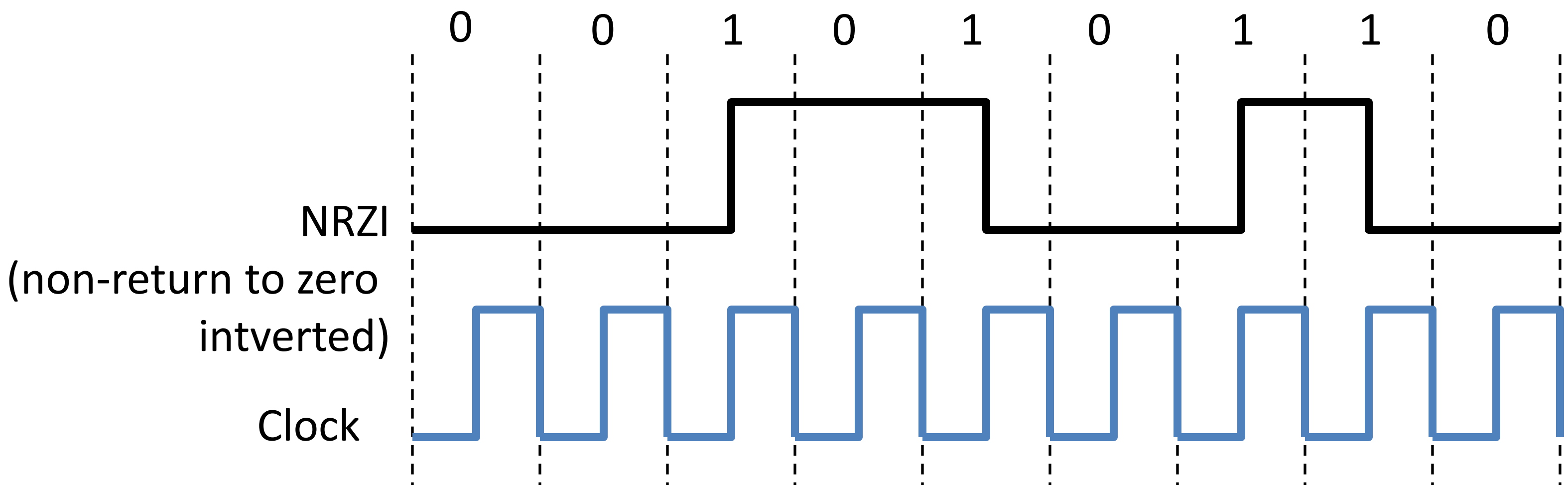


Drawbacks of NRZ

- **No signal could be interpreted as 0 (or vice-versa)**
- **Consecutive 1s or 0s are problematic**
- **Baseline wander problem**
  – How do you set the threshold?
  – Could compare to average, but average may drift
- **Clock recovery problem**
  – For long runs of no change, could miscount periods

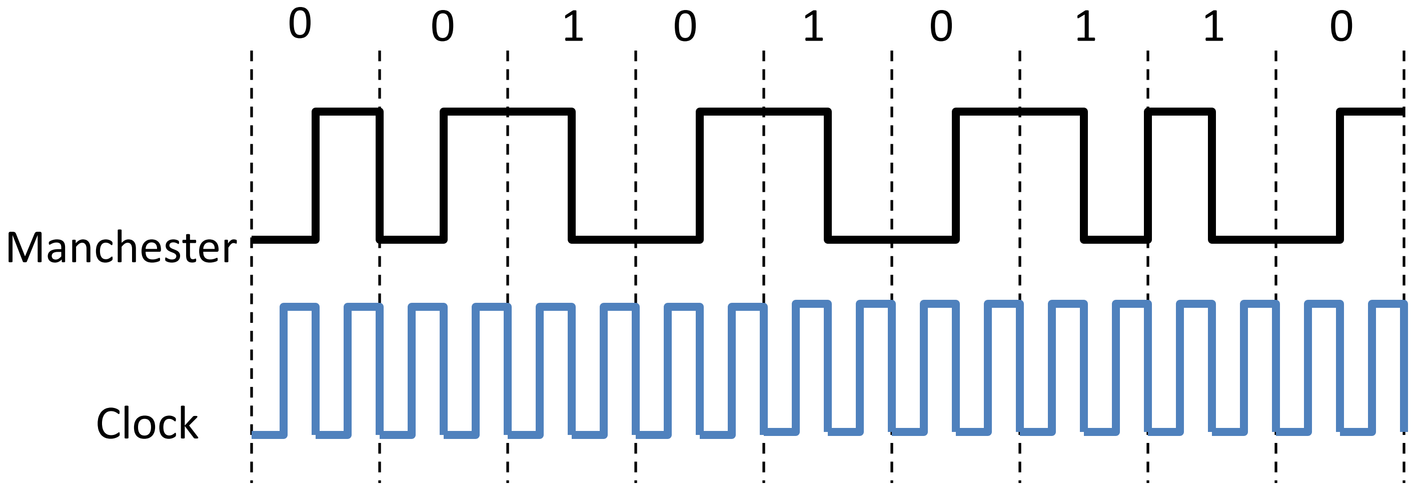
Alternative Encodings

- **Non-return to Zero Inverted (NRZI)**
  - Encode 1 with transition from current signal
  - Encode 0 by staying at the same level
  - At least solve problem of consecutive 1s



Manchester

- **Map 0 → chips 01; 1 → chips 10**
  - Transmission rate now 1 bit per two clock cycles
- **Solves clock recovery, baseline wander**
- **But cuts transmission rate in half**



4B/5B

- **Can we have a more efficient encoding?**
- **Every 4 bits encoded as 5 *chips***
- **Need 16, 5-bit codes:**
  - selected to have no more than one leading 0 and no more than two trailing 0s
  - Never get more than 3 consecutive 0s
- **Transmit chips using NRZI**
- **Other codes used for other purposes**
  - E.g., 11111: line idle; 00100: halt
- **Achieves 80% efficiency**

Encoding

- 4B/5B encoding

0000 → 11110	16 left
0001 → 01001	11111 – when the line is idle
0010 → 10100	00000 – when the line is dead
..	00100 – to mean halt
..	
1111 → 11101	13 left : 7 invalid, 6 for various control signals

Encoding Goals

- **DC Balancing (same number of 0 and 1 chips)**
- **Clock synchronization**
- **Can recover some chip errors**
- **Constrain analog signal patterns to make signal more robust**
- **Want near channel capacity with negligible errors**
  – Shannon says it's possible, doesn't tell us how
  – Codes can get computationally expensive
- **In practice**
  – More complex encoding: fewer bps, more robust
  – Less complex encoding: more bps, less robust

Questions so far?

Photo: Lewis Hine

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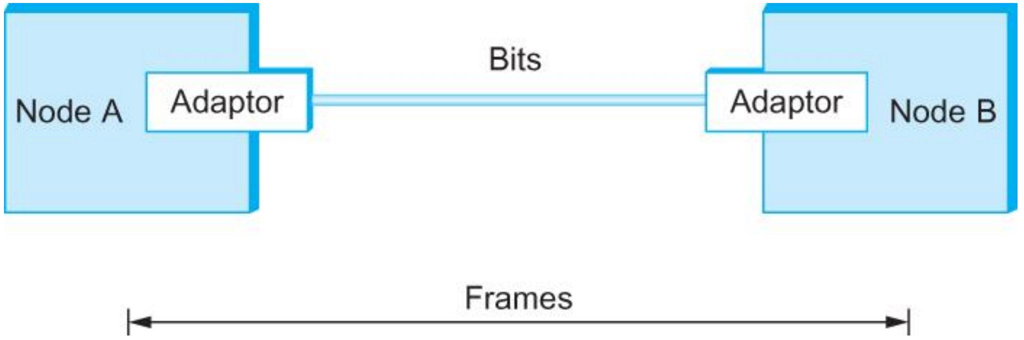
Framing

- Given a stream of bits, how can we represent boundaries?
- Break sequence of bits into a frame
- Typically done by network adaptor



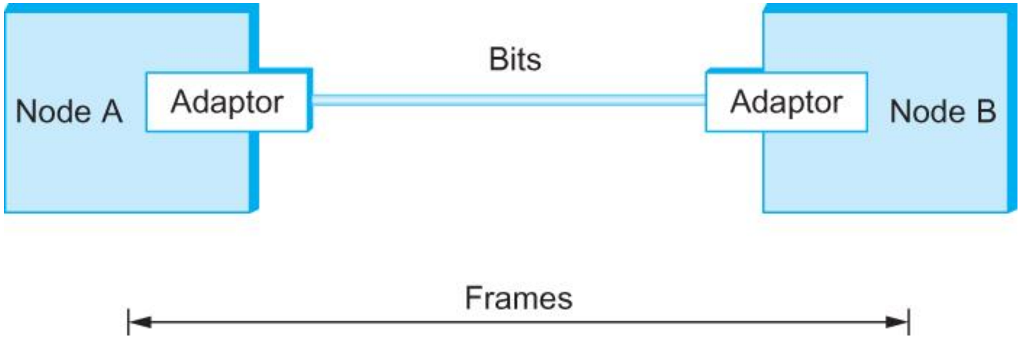
Bits flow between adaptors, frames between hosts

Link Layer

Framing

Representing Boundaries

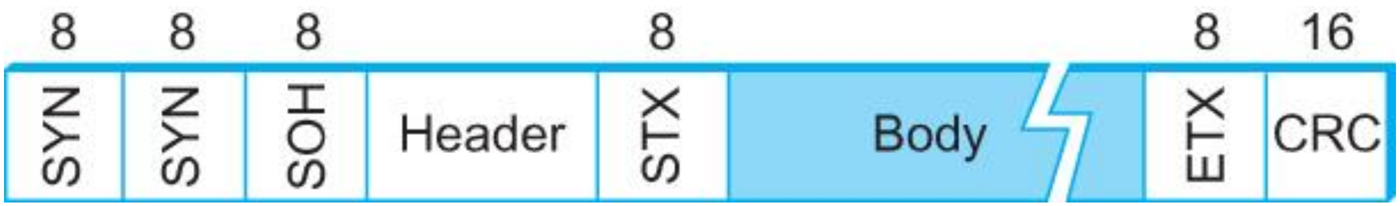
- Sentinels
- Length counts
- Clock-based



Bits flow between adaptors, frames between hosts

Sentinel-based Framing

- **Byte-oriented protocols (*e.g.* BISYNC, PPP)**
  - Place special bytes (SOH, ETX,…) in the beginning, end of messages



- **What if ETX appears in the body?**
  - Escape ETX byte by prefixing DLE byte
  - Escape DEL byte by prefixing DLE byte
  - Technique known as *character stuffing*

Bit-Oriented Protocols

- **View message as a stream of bits, not bytes**
- **Can use sentinel approach as well (*e.g.*, HDLC - High Level Data Link Control)**

  – HDLC begin/end sequence 01111110

- **Use *bit stuffing* to escape 01111110**
  – Always append 0 after five consecutive 1s in data
  – After five 1s, receiver uses next two bits to decide if stuffed, end of frame, or error.

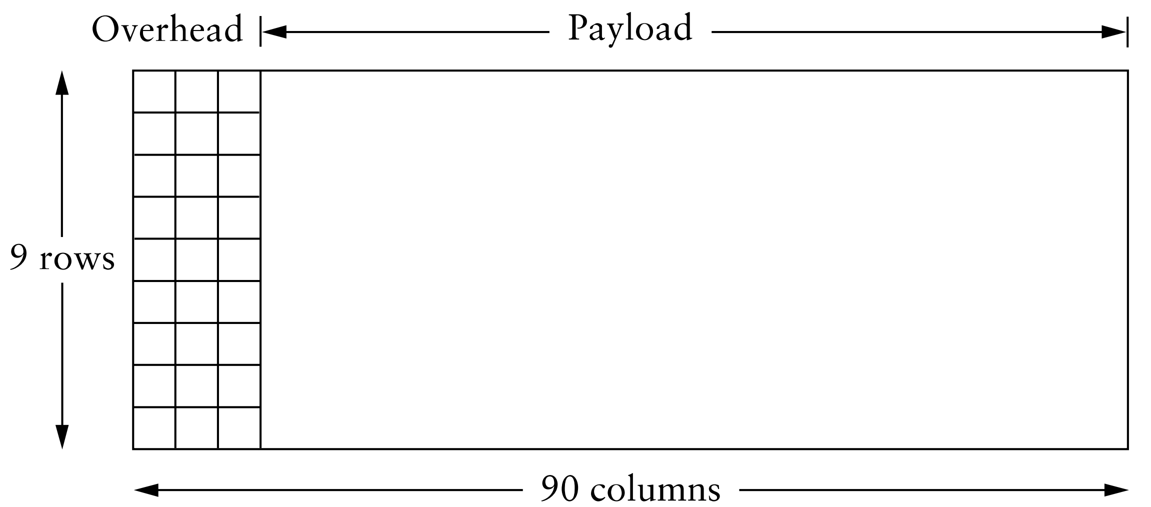
Length-based Framing

- **Drawback of sentinel techniques**
  - Length of frame depends on data
- **Alternative: put length in header (*e.g.*, DDCMP)**

- **Danger: Framing Errors**
  - What if high bit of counter gets corrupted?
  - Adds 8K to length of frame, may lose many frames
  - CRC checksum helps detect error

Clock-based Framing

- ***E.g.*, SONET (Synchronous Optical Network)**
  - Each frame is 125μs long
  - Look for header every 125μs
  - Encode with NRZ, but XOR payload with 127-bit string to ensure lots of transitions



Error Detection

- **Basic idea: use a checksum**
  - Compute small checksum value, like a hash of packet
- **Good checksum algorithms**
  - Want several properties, *e.g.*, detect any single-bit error
  - Details in a later lecture

Next Week

- **Next week: more link layer**
  - Flow Control and Reliability
  - Ethernet
  - Sharing access to a shared medium
  - Switching
- **Thursday Sep 19th: Snowcast due, HW1 out**