

CSCI 1510

- Program Obfuscation (continued)
- Final Review

ANNOUNCEMENT: Course Feedback & Critical Review

Program Obfuscation

Alice



P (program)



Obfuscate



\tilde{P}

```
int E,L,O,R,G[42][m],h[2][42][m],g[3][8],c
[42][42][2],f[42]; char d[42]; void v( int
b,int a,int j){ printf("\33[%d;%df\33[4%d"
"m ",a,b,j); } void u(){ int T,e; n(42)o(
e,m)if(h[0][T][e]-h[1][T][e]){ v(e+4+e,T+2
,h[0][T][e]+1?h[0][T][e]:0); h[1][T][e]=h[
0][T][e]; } fflush(stdout); } void q(int l
,int k,int p){
int T,e,a; L=0
; O=1; while(O
){ n(4&&L){ e=
k+c[l] [T][0];
h[0][L-1+c[l]][
T][1]][p?20-e:
```

Bob



\tilde{P}

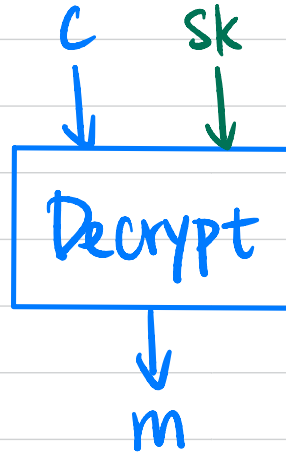
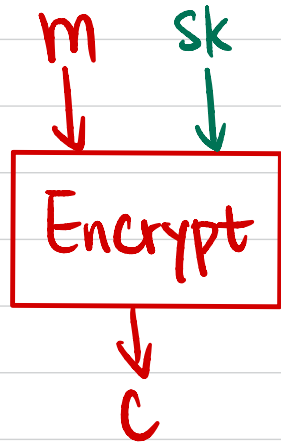


$\tilde{P}(x) \rightarrow y$

$P = ?$

Goal: Make the program "unintelligible" without affecting its functionality.

Symmetric-Key to Public-Key



Formal Definition: Virtual Black Box (VBB)

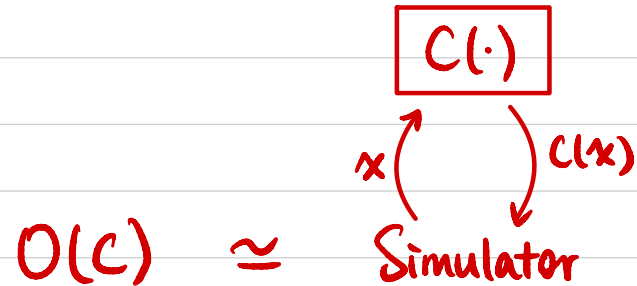
Obfuscator O : $C \xrightarrow{O} O(C)$

• **Functionality:** $O(C)$ computes the same function as C .

• **Polynomial Slowdown:** $|O(C)| \leq \text{poly}(n) \cdot |C|$

• **Security (Virtual Black Box):**

\forall PPT A , \exists PPT S , s.t. $\forall C$, $A(O(C)) \stackrel{c}{\approx} S^{C(\cdot)}(1^{|C|})$.



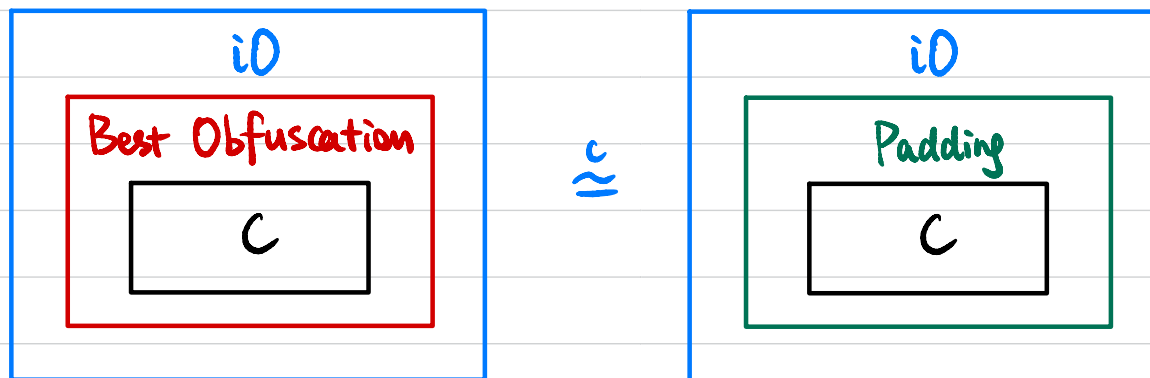
Thm VBB obfuscator for all poly-sized circuits is impossible to achieve.

$$C(x) := \begin{cases} b & \text{if } x=a \\ m & \text{if } x(a)=b \\ 0 & \text{otherwise} \end{cases}$$

Formal Definition: Indistinguishability Obfuscation (iO)

Obfuscator O : $C \xrightarrow{O} O(C)$

- **Functionality**: $O(C)$ computes the same function as C .
- **Polynomial Slowdown**: $|O(C)| \leq \text{poly}(n) \cdot |C|$
- **Security (indistinguishability obfuscation)**:
If C_0 & C_1 compute the same function and $|C_0| = |C_1|$,
then $O(C_0) \stackrel{c}{\approx} O(C_1)$
- **Best Possible Obfuscation**



PKE from iO

Let $G: \{0,1\}^n \rightarrow \{0,1\}^{2n}$ be a length-doubling PRG.

• $\text{Gen}(1^n)$:

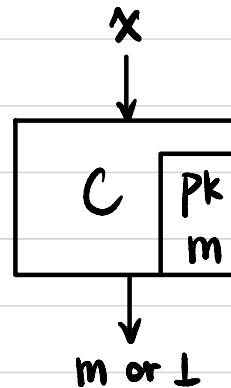
$$sk \leftarrow \{0,1\}^n$$

$$pk := G(sk)$$

• $\text{Enc}_{pk}(m)$:

$$C_{pk,m}(x) := \begin{cases} m & \text{if } G(x) = pk \\ \perp & \text{otherwise} \end{cases}$$

$$\text{Output } c \leftarrow \text{iO}(C_{pk,m})$$



• $\text{Dec}_{sk}(c)$: ?

Thm If G is a PRG and $\text{iO}(\cdot)$ is an indistinguishability obfuscator, then this PKE scheme is CPA-secure.

Is it possible?

- 2001: Notion introduced
- 2013: First "candidate" construction from multilinear maps
- 2013-2020: Attack, fixes, new constructions from new assumptions
- 2020: New construction from well-founded assumptions

Final Review

- Cryptographic Hardness Assumptions
 - Factoring / RSA Assumptions
 - DLOG / CDH / DDH Assumptions
 - LWE Assumption (Post-Quantum)
- Key Exchange
 - Definition
 - Construction: Diffie-Hellman
- Public-Key Encryption
 - Definition: CPA / CCA
 - Constructions: El Gamal / RSA / Regev

Final Review

- Theoretical Assumptions
 - One-Way Function / Permutation: Definition & Candidates
 - Hard-Core Predicate: Definition & Construction
 - PRG / PRF from OWP
 - Trapdoor Permutation: Definition & Candidate (RSA)
 - PKE from TDP
- Fully Homomorphic Encryption
 - Definition & Applications
 - Somewhat Homomorphic Encryption over Integers & from LWE (GSW)
 - Bootstrapping SWHE to FHE

Final Review

- Digital Signature
 - Definition
 - Hash-and-Sign Paradigm
 - Construction 1: RSA-FDH
 - Proof in the Random Oracle Model
 - Construction 2: Schnorr
 - Identification Scheme: Definition & Construction from DLOG (Schnorr)
 - Fiat-Shamir Transform

Final Review

- Zero-Knowledge Proof
 - Definition: Completeness / Soundness / Zero-Knowledge
 - Example: ZKP for Diffie-Hellman Tuples
 - Proof Technique: Rewinding
 - ZKP for All NP (Graph 3-Coloring)
 - Commitment Scheme
 - Non-Interactive ZK

Final Review

- Secure Multi-Party Computation
 - Definition: Semi-Honest / Malicious
 - Applications
 - Example: Private Set Intersection from DDH
 - MPC for Any Function (GMW)
 - Oblivious Transfer: Definition & Construction from CDH
- Program Obfuscation
 - Definitions: VBB / iO
 - Example: PKE from iO

THANK YOU 😊