# **Lecture 4: Synchronization Primitives**

# CS178: Programming Parallel and Distributed Systems

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I. Overview

## A. Topics covered last time

- 1. Different types of synchronization to worry about
- 2. Some sense as to why this is difficult
- **B.** Topics to cover today
  - 1. Primitives used for multithreaded programming
  - 2. Java multithreaded programming
  - 3. Start talking about design and implementation techniques
- C. For those who see much of this as a review consider the problem of using multiple threads to implement a prime number sieve
  - 1. How would you make maximal use of threads
  - 2. What are the synchronization problems

## **II. Threads**

- A. A thread is a virtual CPU executing code in the process
  - 1. Has its own set of registers
  - 2. Has its own local variables / stack
  - 3. Has its own location counter

## **B.** Basic thread operations

- 1. Create
  - a) Sets up a new thread
  - b) Creates its stack
  - c) Identifies what code it should execute
- 2. Start

a) Actually starts the thread executing

#### 3. Get current thread

a) Returns handle to the current executing thread

## 4. Associate data with the thread

- a) If object-based, this is in the thread object
- b) Otherwise, system provides an associative table of thread-specific information

#### 5. Exit

- a) Ends the thread with a value
- b) Equivalent to exiting the top-level routine of the thread

### 6. WaitFor

- a) Waits for one or more threads to terminate
- b) Used to synchronize threads at end of operation

### 7. Safe interrupt or stop

- a) Means for having one thread control another
- b) Inherently unsafe
- c) Thus this is typically done via notification and polling

## C. C++ threads

## 1. Pthreads library -- C-style interface to threads

## 2. Class interface

- a) MS CWinThread class
- b) Unix -- /pro/bloom/stdlib/src/bloom\_pthread.H

## **D.** Java threads

- 1. java.lang.thread class is standard
- 2. Methods provided for all the above operations
- 3. User defines a new subclass inheriting from this
  - a) Subclass contains a "run" method that is the body
  - b) Subclass can define thread-local storage as appropriate

# **III.Synchronization Primitives**

# A. Problem

- 1. Memory synchronization
- 2. Preventing threads from interfering with each other

#### 3. Variety of solutions

- a) All Turing equivalent
- b) All provide for non-busy wait in the OS
- c) All let the operating system ensure fairness

## **B.** Mutex

1. Basic uses: Critical region

### 2. Operations

- a) Lock
- b) Unlock
- c) (test lock) -- lock without the wait

## **C.** Semaphore: mutex with a counter

### 1. Uses

- a) Can be used for a critical region
- b) Can be used to control a queue or stack as well

### 2. Operations

- a) V -- increment the counter (unlock)
- b) P -- decrement the counter if > 0, else wait (lock)

## **D. RWLock: read/write lock**

#### 1. Uses

a) Can be used where writes block readers, readers don't block each other, any reader blocks a writer

#### 2. Operations

- a) readlock, writelock
- b) unlock

## 3. Implementing RW locks using semaphore/mutex

- a) Keep a global variable of # of readers
- b) Reader entry -- as a critical region, update this counter, if it was 0, set the writer lock
- c) Reader exit -- as a critical region, decrement the counter, if it is now 0, unset the writer lock

# E. Condition variables

#### 1. Uses

- a) Sometimes we want to check a condition in a critical region and then wait for it to be valid
- b) But we don't want to wait in the critical region
- c) Rather than exiting (and then the check becoming invalid) or busy looping, we use condition variables

#### 2. Operations

- a) wait(mutex)
- b) signal -- wake up the next thread waiting on the cv
- c) broadcast -- wake up all threads waiting on the cv
- d) Note that this handles the mutex correctly

## F. Monitors

#### 1. The above are all data structures, not program structures

- a) Monitors are a language construct that embodies these
- b) Essentially a monitor is a block that is protected by a mutex and that supports condition variables
- c) Block provides a program scope, local variables (that are shared), multiple entry points, ...
- 2. Used in modula (& modula 2, 3) & its derivatives
- 3. Sort of used in Java

## G. Java Synchronization

#### 1. Every java object is a potential mutex

- a) Ability to define critical regions for the object
- b) Either from synchronized methods of the object
- c) Or by using synchronized(object) { ... }

#### 2. Every java object is a potential condition variable

- a) wait() or wait(timeout) is the wait call
  - (1) Must be done in a synchronized context
- b) notify() and notifyAll() represent signal and broadcast

## 3. Every java object is a potential monitor

- a) Can provide a set of mutually exclusive entry points
- b) Local storage is shared

- c) Language construct
- d) Supports condition variables

# **IV. Multithreaded Java Programming**

## A. Basic concepts

#### 1. Determine what can be done in parallel

a) What is the task of each thread

#### 2. Determine what are the shared data structures

- a) Attempt to minimize sharing
- b) Fewest possible shared data structures
- c) Least amount of sharing within the data structure

#### 3. Determine how to create/stop the threads

- a) Directly as needed (with waitfor at end)
- b) Using a queue of tasks

## **B.** Problems to be aware of

#### 1. Deadlock

- a) Dining philosophers problem
- b) Recursive data structures

#### 2. Performance

- a) Synchronization is not free
- b) What else needs to be synchronized (beyond threads)
  - (1) Memory management
  - (2) I/O

## C. Example

#### 1. Consider the problem of implementing a prime number sieve

- a) Generate list of prime numbers
- b) For each number, check it against each number in the list up to its square root

#### 2. What can be done in parallel

- 3. What are the shared data structures
- 4. How to create/stop threads