

CS2: A Searchable Cryptographic Cloud Storage System

Seny Kamara (MSR)
Charalampos Papamanthou (UC Berkeley)
Tom Roeder (MSR)

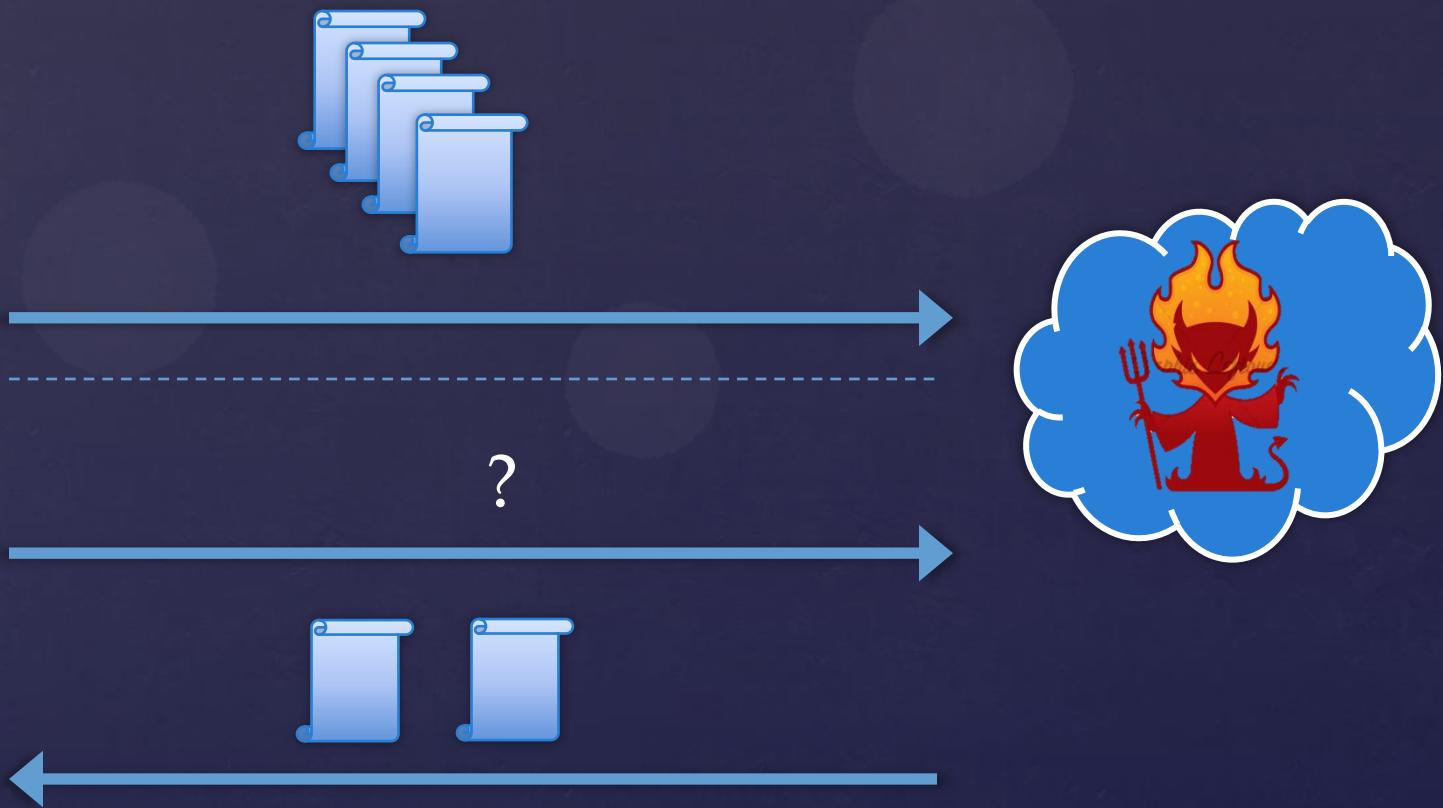
Cloud Computing



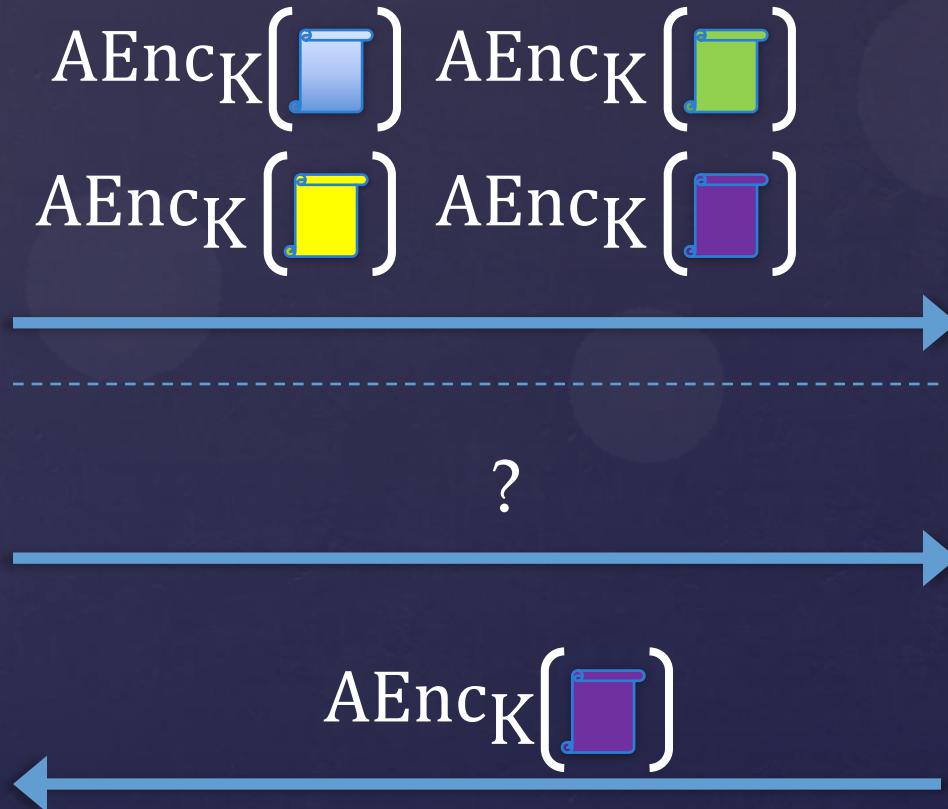
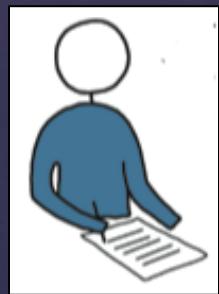
Cloud Computing

- Main concern
 - *will my data be safe?*
 - will anyone see it?
 - can anyone modify it?
- Security solutions
 - VM isolation
 - Single-tenant servers
 - Access control
 - ...
- Cloud provides *stronger* security than self-hosting [Molnar-Schecter-10]
- Q: but what if I don't trust the *cloud operator*?

Cloud Storage



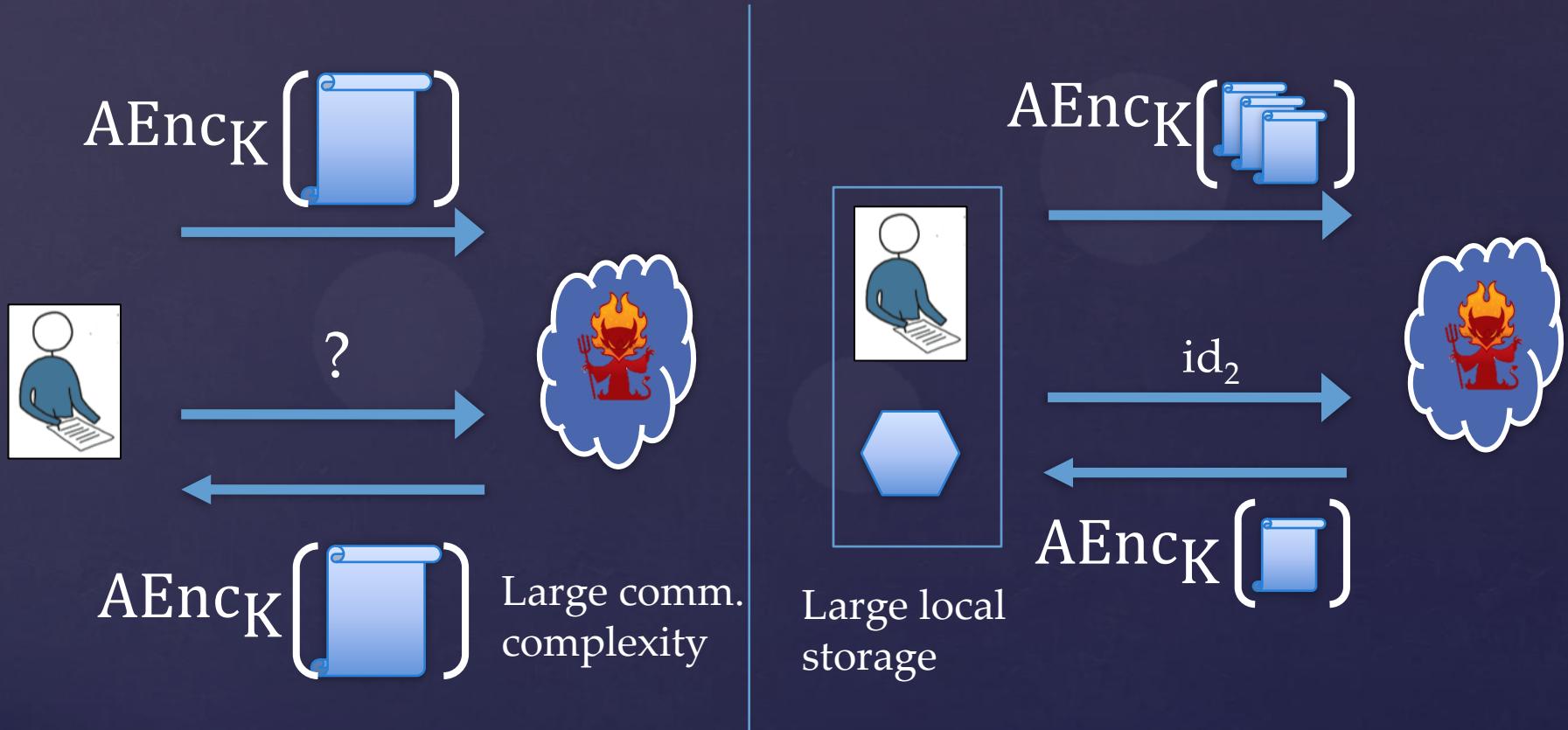
Traditional Approach



Search-based Access

- File-based access is hard (esp. for large data)
- Search-based access is preferred
 - Web search
 - Desktop search
 - Apple Spotlight, Google Desktop, Windows Desktop
 - Enterprise search

Two Simple Solutions to Search



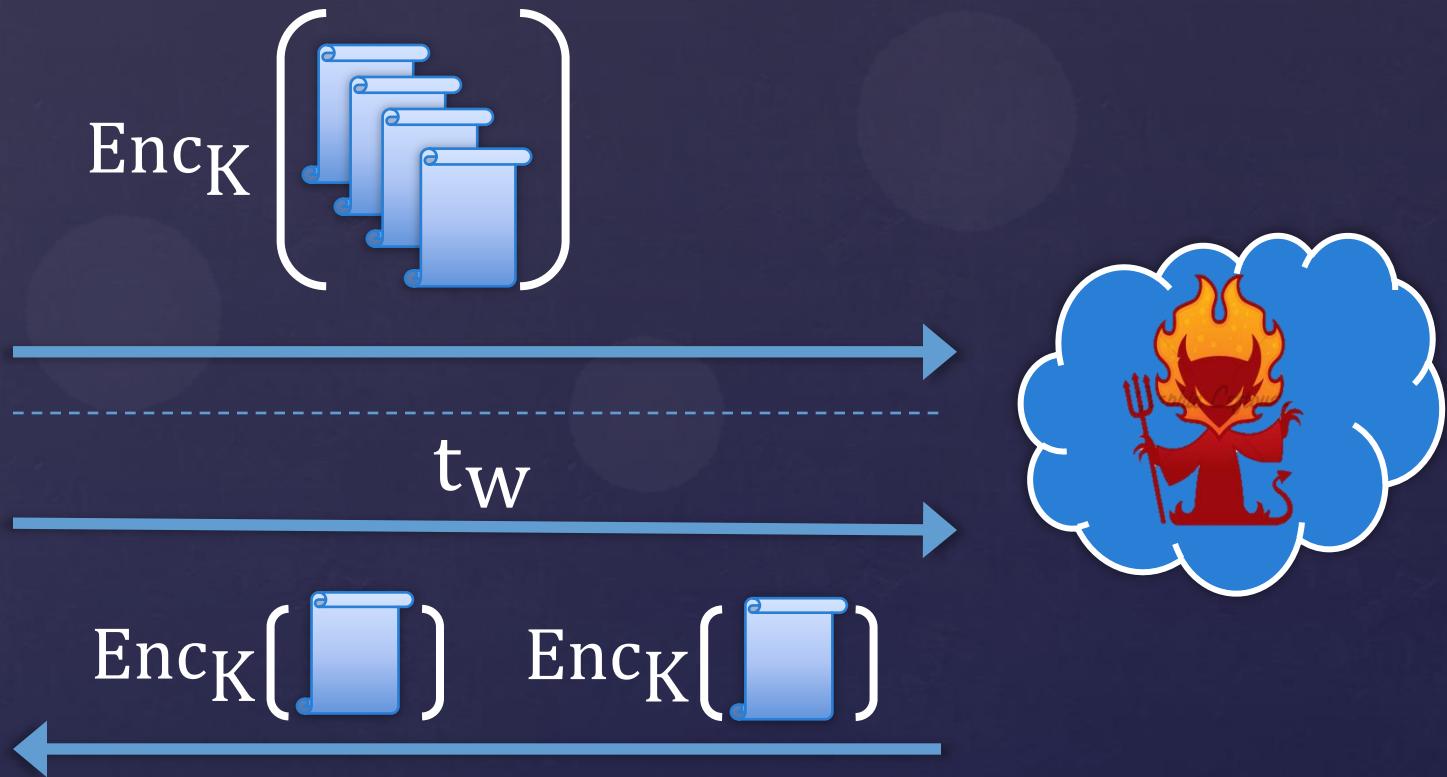
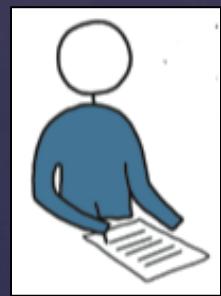
Q: can we achieve the best of both?

Outline

- **Motivation**
- **CS2 building blocks**
 - Symmetric searchable encryption
 - Search authenticators
 - Proofs of storage
- **CS2 Protocols**
 - for standard search
 - for assisted search
- **Experiments**

CS2 Building Blocks

Searchable Symmetric Encryption [SWP01]

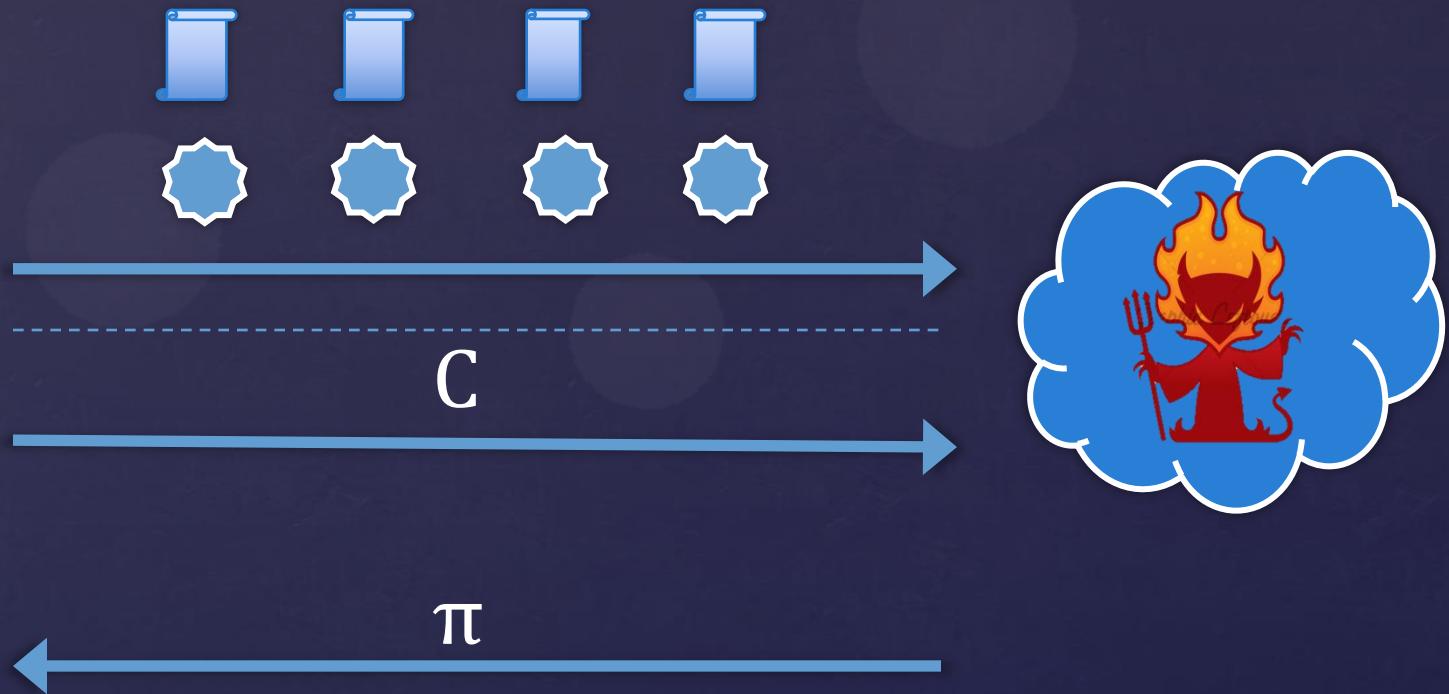
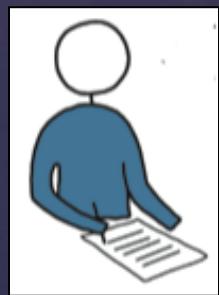


Searchable Symmetric Encryption

- [Goldreich-Ostrovsky-96]
 - ☺: hides everything
 - ☹: interactive
- [Song-Wagner-Perrig-01]
 - ☺: non-interactive
 - ☹: static, *linear search time, leaks information*
- [Goh03, Chang-Mitzenmacher-05]
 - ☺: non-interactive, dynamic
 - ☹: *linear search time, non-adaptive security* (CKA1-security)
- [Curtmola-Garay-K-Ostrovsky-06]
 - ☺: non-interactive, sub-linear search (optimal), adaptive security
 - ☹: *static*

We need new SSE!

Proofs of Storage [ABC+07, JK07]

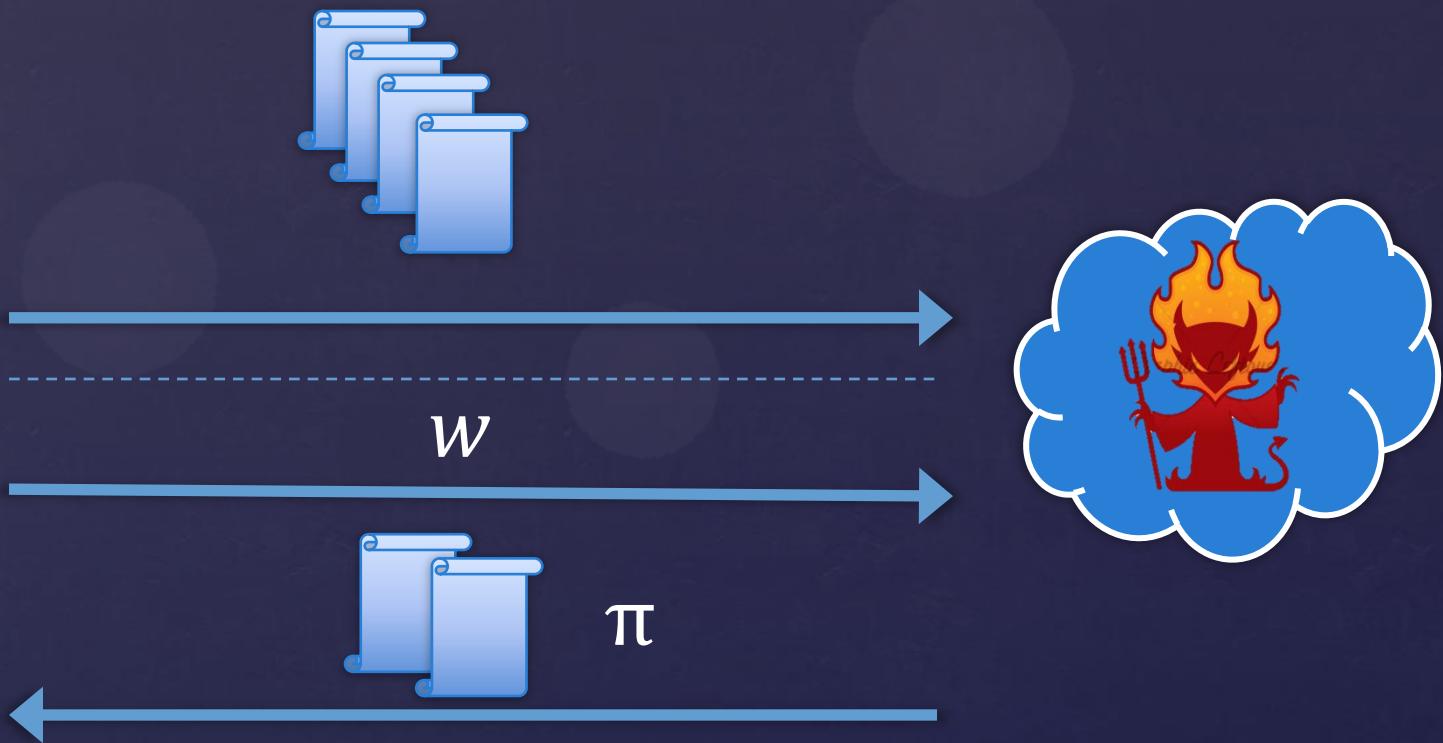
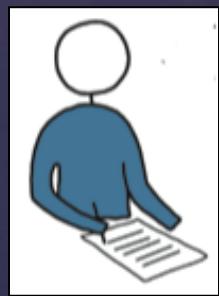


Proofs of Storage

- [ABC+07,JK07,SW08,DVW09,AKK09]
 - ☺: efficient
 - ☹: static
- [APMT08]
 - ☺: efficient and dynamic
 - ☹: bounded verifications
- [EKPT09]
 - ☺: efficient, dynamic, unlimited verification
 - ☹: patented

We need new PoS!

Search Authenticator



Search Authenticators

- [GGP10, CVK10, CVK11]
 - ☺: general-purpose
 - ☹: inefficient (due to FHE) & static
- [CRR11]
 - ☺: general-purpose, efficient
 - ☹: requires two non-colluding clouds
- [BGV11]
 - ☹: proof generation is linear & static

We need new VC/SA!

Outline

- **Motivation**
- **CS2 building blocks**
 - *Symmetric searchable encryption*
 - Search authenticators
 - Proofs of storage
- **CS2 Protocols**
 - for standard search
 - for assisted search
- **Experiments**

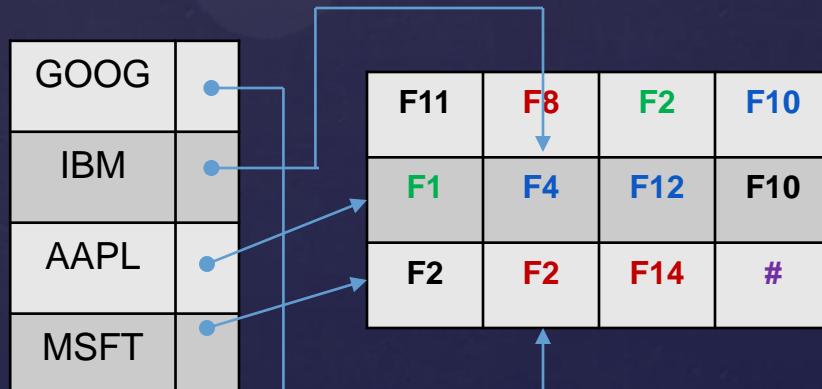
SSE-1 [CGKO06]

MSFT	F2	F10	F11
GOOG	F2	F8	F14
AAPL	F1	F2	
IBM	F4	F10	F12

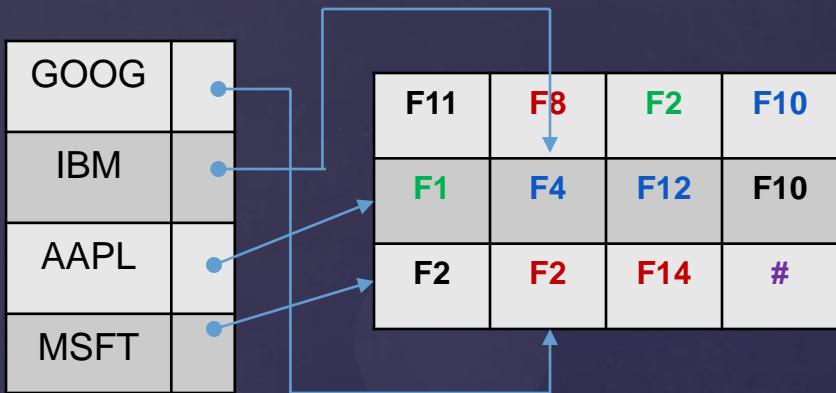
1. Build inverted/reverse index

Posting list

2. Randomly permute array & nodes

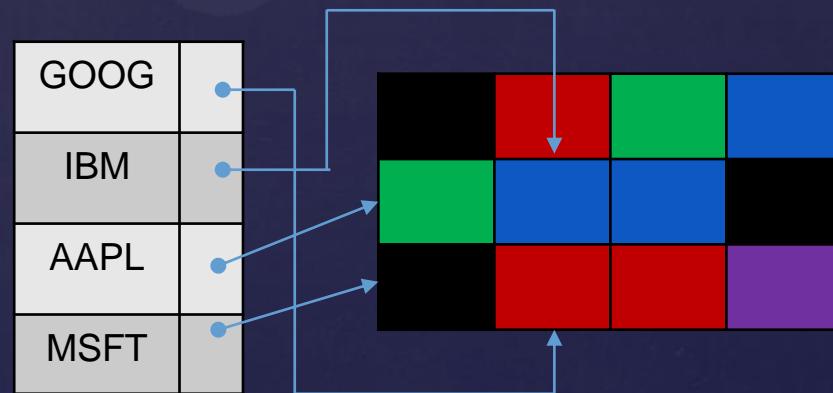


SSE-1 [CGKO06]

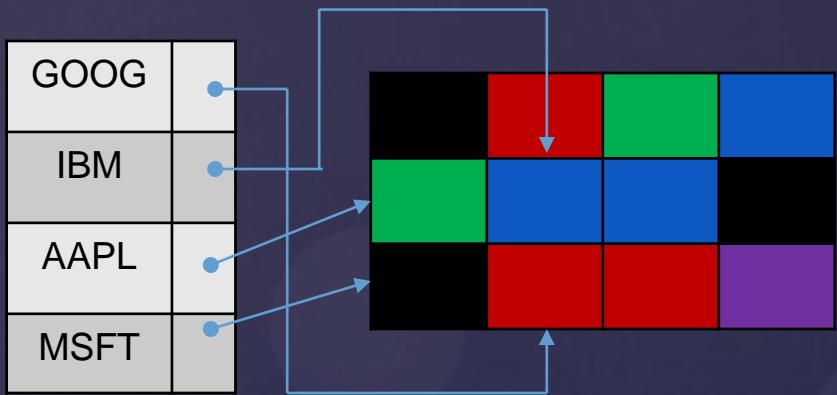


2. Randomly permute array & nodes

3. Encrypt nodes



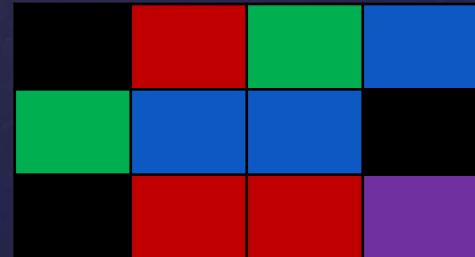
SSE-1 [CGKO06]



3. Encrypt nodes

4. "Hash" keyword & encrypt pointer

$F_K(GOOG)$	$Enc(\bullet)$
$F_K(IBM)$	$Enc(\bullet)$
$F_K(AAPL)$	$Enc(\bullet)$
$F_K(MSFT)$	$Enc(\bullet)$

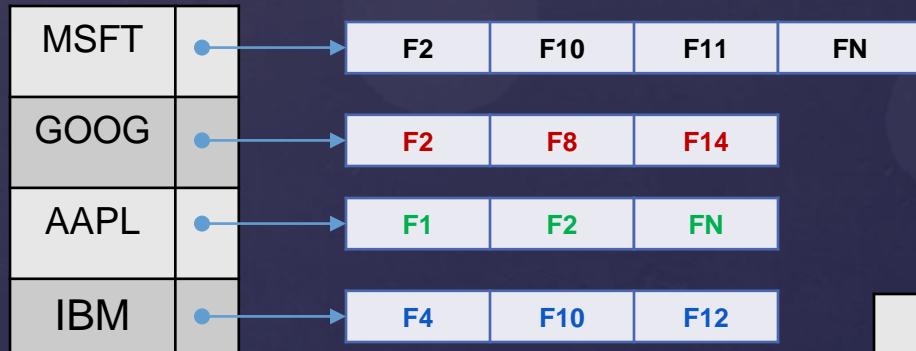


Limitations of SSE-1

- Non-adaptively secure \Rightarrow adaptive security
 - Idea #1 [Chase-K-10]
 - replace encryption scheme with symmetric non-committing encryption
 - only requires a PRF + XOR
 - ☺: doesn't work for dynamic data
 - Idea #2
 - Use RO + XOR

Limitations of SSE-1

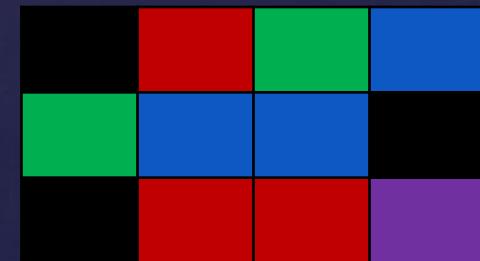
- Static data \Rightarrow dynamic data
- Problem #1:
 - given new file $F_N = (\text{AAPL}, \dots, \text{MSFT})$
 - append node for F to list of every w_i in F



1. Over unencrypted index

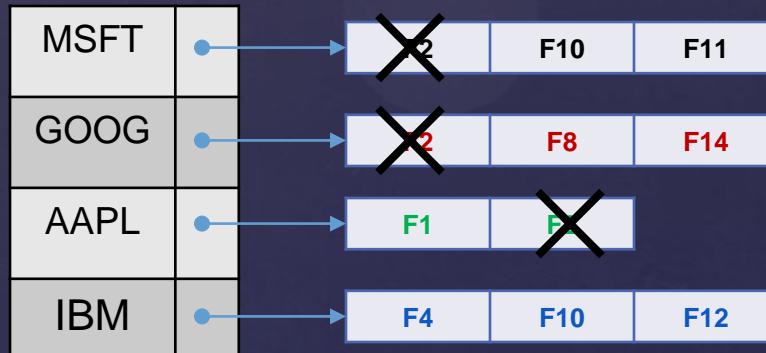
2. Over encrypted index ???

$F_K(\text{GOOG})$	$\text{Enc}(\bullet)$
$F_K(\text{IBM})$	$\text{Enc}(\bullet)$
$F_K(\text{AAPL})$	$\text{Enc}(\bullet)$
$F_K(\text{MSFT})$	$\text{Enc}(\bullet)$



Limitations of SSE-1

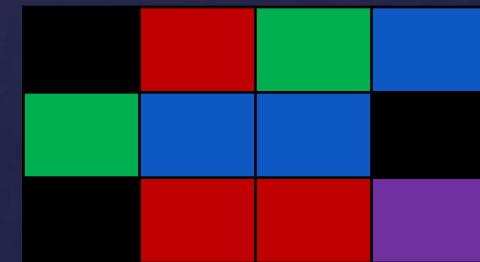
- Static data \Rightarrow dynamic data
- Problem #2:
 - When deleting a file $F_2 = (\text{AAPL}, \dots, \text{MSFT})$
 - delete all nodes for F_2 in every list



1. Over unencrypted index

2. Over encrypted index ???

$F_K(\text{GOOG})$	$\text{Enc}(\bullet)$
$F_K(\text{IBM})$	$\text{Enc}(\bullet)$
$F_K(\text{AAPL})$	$\text{Enc}(\bullet)$
$F_K(\text{MSFT})$	$\text{Enc}(\bullet)$



Limitations of SSE-1

- Static data \Rightarrow dynamic data
 - Idea #1
 - Memory management over encrypted data
 - Encrypted free list
 - Idea #2
 - List manipulation over encrypted data
 - Use homomorphic encryption (here just XOR) so that pointers can be updated obliviously
 - Idea #3
 - deletion is handled using an “dual” SSE scheme
 - given deletion/search token for F_2 , returns pointers to F_2 ’s nodes
 - then add them to the free list homomorphically

Outline

- **Motivation**
- **Related work & our approach**
- **CS2 building blocks**
 - Symmetric searchable encryption
 - *Search authenticators*
 - Proofs of storage
- **CS2 Protocols**
 - for standard search
 - for assisted search
- **Experiments**

Limitations of Verifiable Computation

- Inefficient \Rightarrow practical
 - Idea #1
 - Design special-purpose scheme (i.e., just for verifying search)
 - Idea #2
 - Use Merkle Tree “on top” of inverted index
 - For keyword w : we efficiently verify its posting list and associated files
 - Generating proof is $O(w^*)$ instead of $O(n)$
- Static \Rightarrow dynamic
 - Idea #1
 - Replace bottom hash with *incremental* hash
 - [Bellare-Goldreich-Goldwasser94, Bellare-Micciancio97]

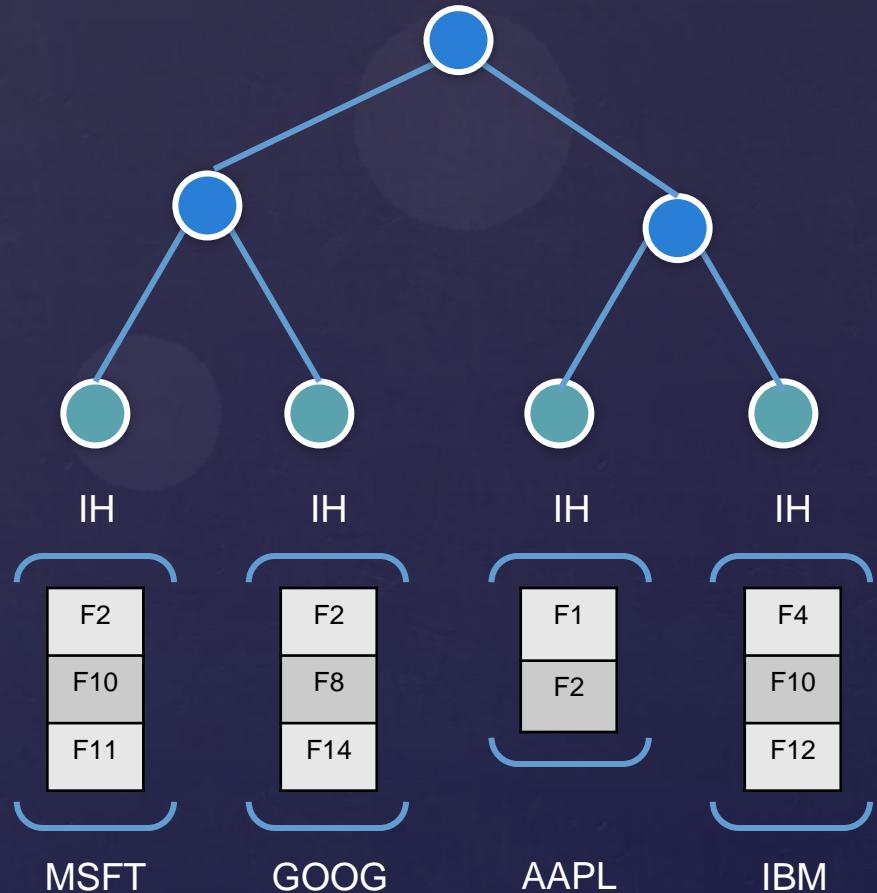
Search Authenticators

1. Build inverted/reverse index

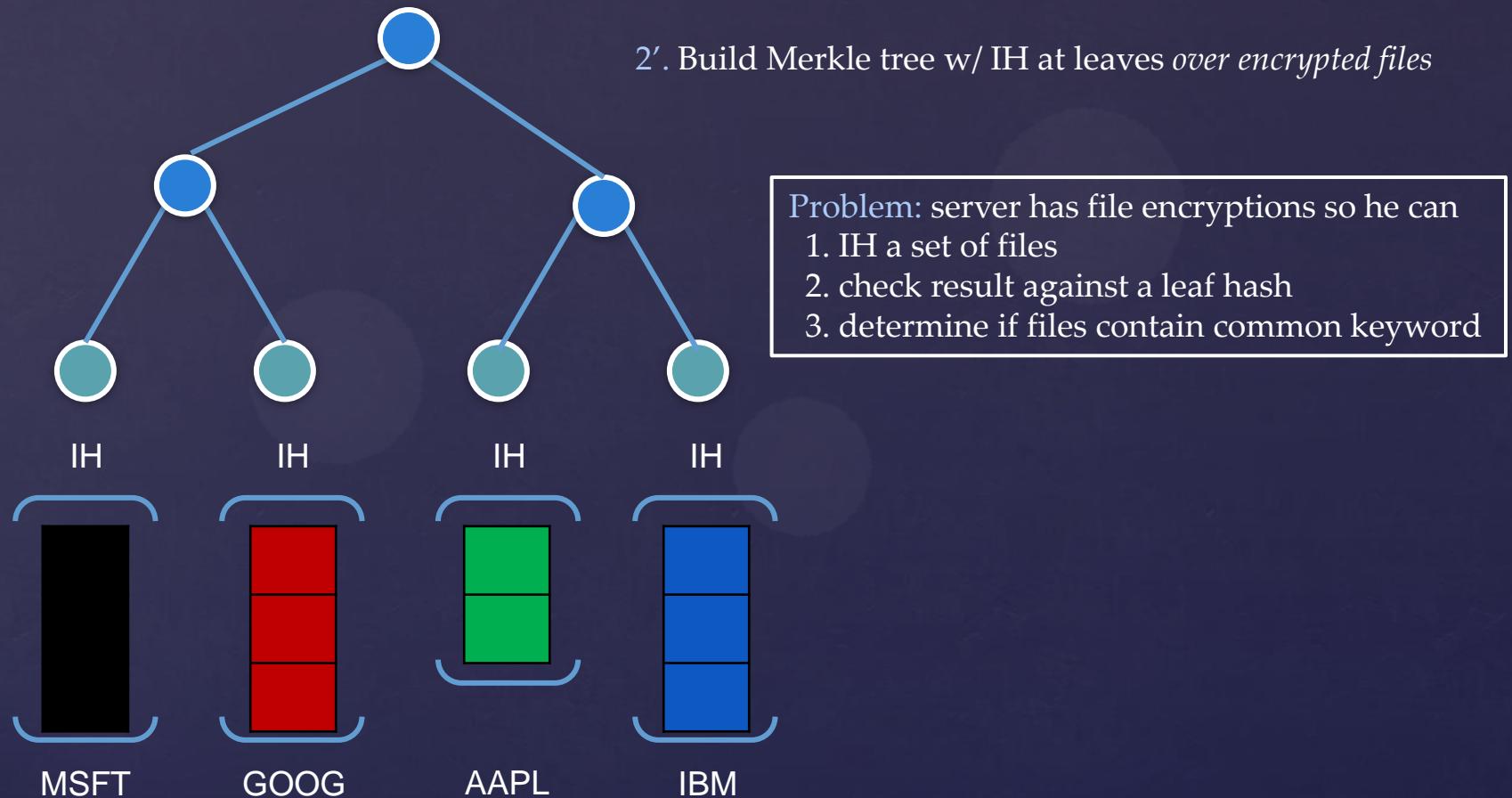
MSFT	F2	F10	F11
GOOG	F2	F8	F14
AAPL	F1	F2	
IBM	F4	F10	F12

2. Build Merkle tree w/ IH at leaves

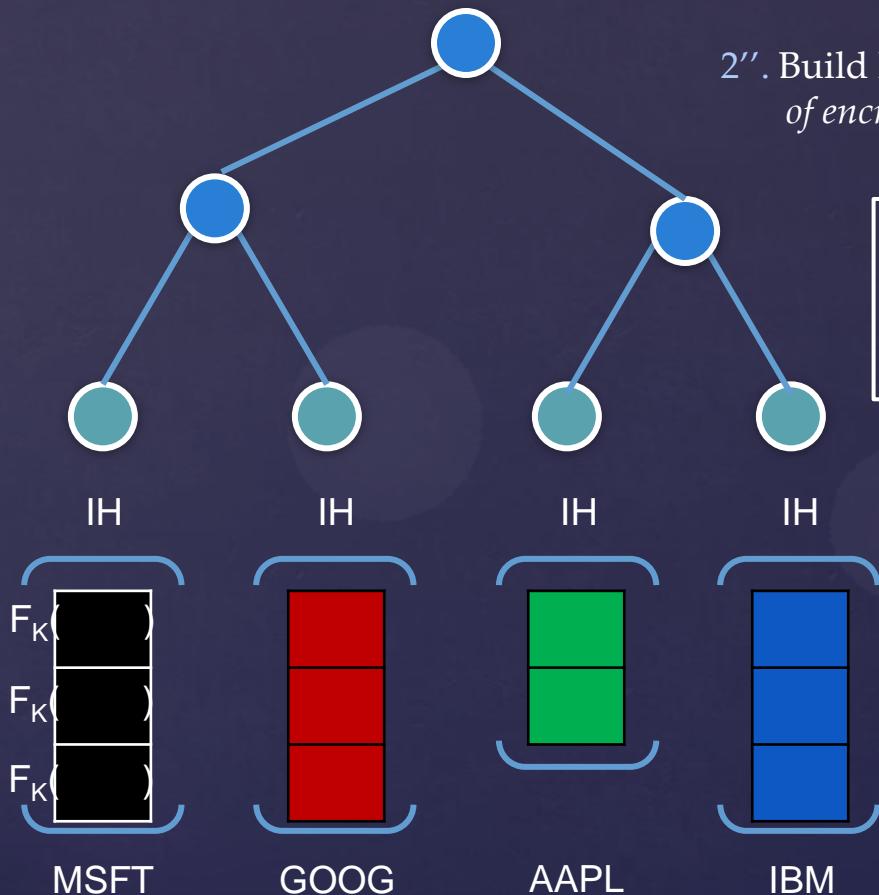
Problem: hash functions are not hiding!



Search Authenticators



Search Authenticators



2''. Build Merkle tree w/ IH at leaves *over keyed hash of encrypted files*

Problem: server has file encryptions so he can

1. IH a set of files
2. check result against a leaf hash
3. determine if files contain common keyword

Proofs of Storage

CS2 Protocols

CS2 Protocols

- **Standard search**
 - User searches for w
 - Server returns documents w/ w
 - Relatively straightforward combination of (dynamic) SSE, PoS & SA
- **Assisted search**
 - User searches for w
 - Server returns summaries of files with w
 - User chooses a subset to retrieve
 - Server returns subset of files with w
 - More complex combination of (dynamic) SSE, PoS, SA + CRHF
 - *Search can be more efficient* (since less data is returned)

CS2 Protocols

- Definitions in ideal/real-world model
 - Cloud storage w/ standard search
 - Cloud storage w/ assisted search
 - ☺
 - easier to use within larger protocols (i.e., *hybrid security models*)
 - Single definition for all desired properties
 - guarantees composition of underlying primitives is OK
 - ☹: definitions & proofs are complicated
- Protocols make black-box use of primitives
 - ☺: modularity -- replace underlying primitives

Experiments

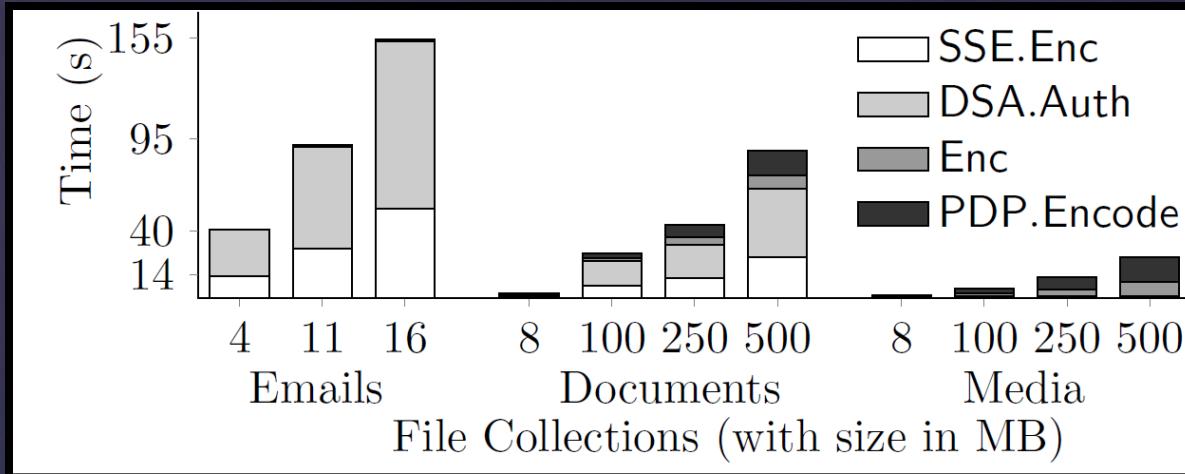
Implementation

- C++
- Microsoft Cryptography API: Next Generation
 - RO: SHA256
 - PRFs: HMAC-SHA256
 - SKE: 128-bit AES/CBC
- Bignum library
 - Prime fields
- We test only the crypto overhead
 - No file transfers over network
 - No reading from disk
 - No indexing costs

Experiments

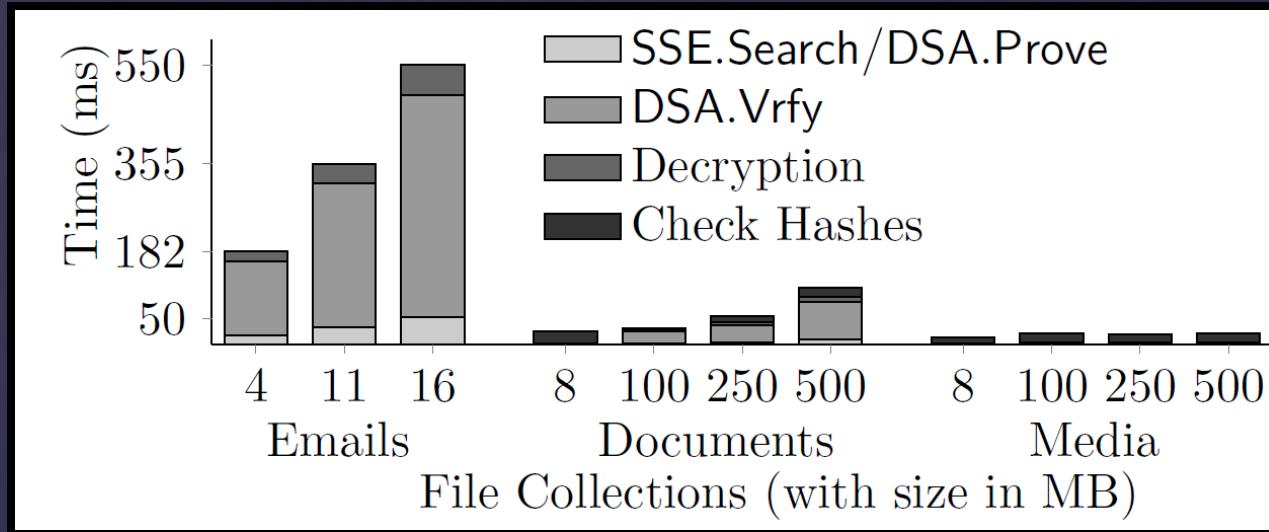
- Intel Xeon CPU 2.26 GHz
 - Windows Server 2008
- 4 datasets
 - Email (enron): 4MB, 11MB, 16MB
 - ≈ every byte is a word
 - Office docs: 8MB, 100MB, 250MB, 500MB
 - Relatively few keywords
 - Media (MP3,WMA, JPG,...): 8MB, 100MB, 250MB, 500MB
 - Barely any keywords
- Average over 10 executions

STORE



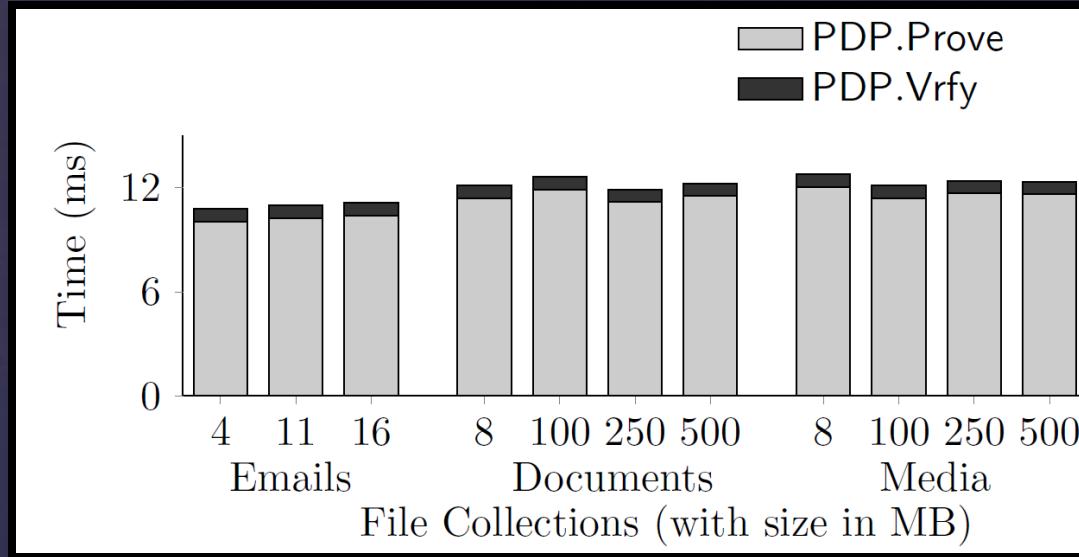
- Total
 - Email (16MB): 2 mins
 - Office (500MB) :1.5 mins
 - Media (500MB): 30 s
 - Email (16GB): 40/15 hours
- Distribution
 - Verifiability: 2/3 of cost
 - SSE: 1/3 cost
 - PoS: negl

SEARCH



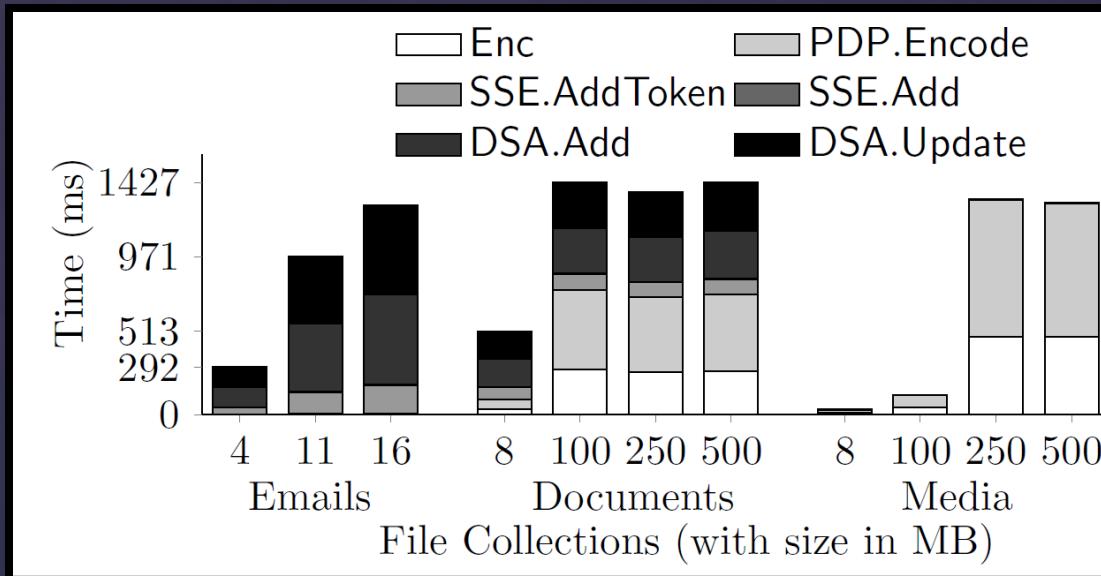
- Total
 - Email (16MB): 0.5 secs
 - Office (500MB): 0.1 secs
 - Media (500MB): 0.025 secs
- Distribution
 - Client verification: 80%
 - Client decryption: 10%
 - Server search + proof: 10%

CHECK



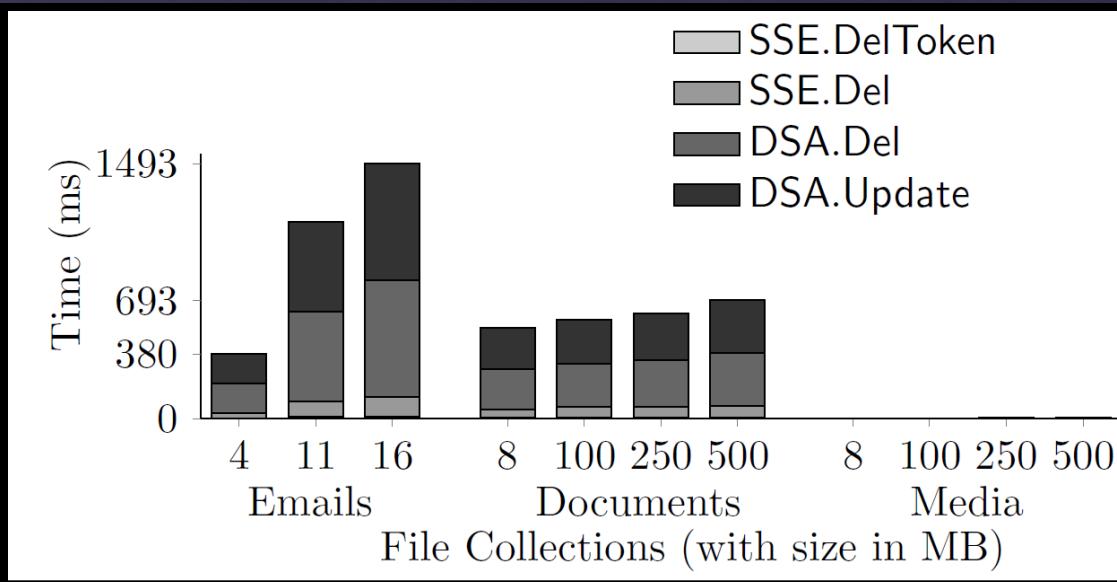
- Total
 - Email (16MB): 12 secs
 - Office (500MB): 12 secs
 - Media (500MB): 12 secs
- Distribution
 - Server Proof: 95%
 - Client verify: 5%

ADD



- Total
 - Email (16MB): 1.5 secs
 - Office (500MB): 1.5 secs
 - Media (500MB): 1.5 secs
- Distribution
 - Email (16MB)
 - 40% client auth state update
 - 40% server auth update
 - 20% add token

DELETE



- Total
 - Email (16MB): 1.5 secs
 - Office (500MB): 0.7 secs
 - Media (500MB): negl
- Distribution
 - 40% server auth update
 - 40% client auth update
 - 20% server index update

Summary

- New Crypto
 - Dynamic and CKA2-secure SSE with sub-linear search
 - Sub-linear verifiable computation for search
 - Unbounded dynamic PDP
- New Protocols
 - Ideal/real-world definitions for secure cloud storage
 - Protocol for standard search
 - Protocol for assisted search
- Implementation & experiments
 - First experimental results for sub-linear SSE
 - Identified verification as bottleneck
 - Office docs seem to be the best workload

Questions?